

$$\begin{bmatrix} x[0] \\ x[1] \\ \vdots \\ x[N-1] \end{bmatrix} = W_N^{-1} \begin{bmatrix} X[0] \\ X[1] \\ \vdots \\ X[N-1] \end{bmatrix}$$

$$W_N^{-1} = \frac{1}{N} W_N^* \leftarrow \text{complex conjugate transpose}$$

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi k}{N} n}, \quad n=0, 1, \dots, N-1.$$

(Inv. D.T.F.T: $x[n] = \frac{1}{2\pi} \int_{-\pi}^{\pi} X(e^{j\omega}) e^{j\omega n} d\omega$)

* What happens if n is outside the index set $0, 1, \dots, N-1$

Define: $x[N] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi k}{N} N} = \frac{1}{N} \sum_{k=0}^{N-1} X[k]$
 $x[0] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi k}{N} \cdot 0} = \frac{1}{N} \sum_{k=0}^{N-1} X[k]$

Define $x[N+1] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi k}{N} (N+1)} = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi k}{N}} = x[1]$

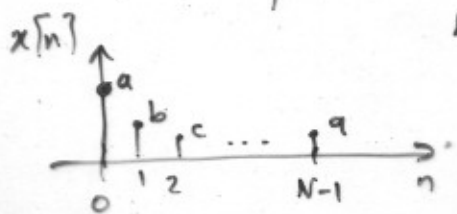
Define $x[-1] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi k}{N} (-1)} = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{-j \frac{2\pi k}{N}} = x[N-1]$

\Rightarrow periodic extension: $x_p[n]$

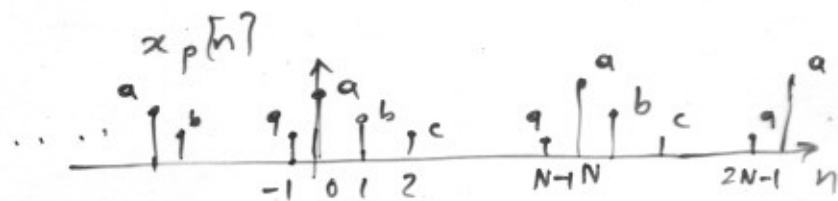
$$x_p[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] e^{j \frac{2\pi k}{N} n}, \quad n=0, \pm 1, \pm 2, \dots$$

$x_p[n] = x[n], \quad \text{for } n=0, 1, 2, \dots, N-1.$

$x_p[n] = \sum_{l=-\infty}^{\infty} x[n - lN], \quad (\text{In general})$



\longrightarrow



FFT for $N=3p$

$$X[k] = \sum_{n=0}^{N-1} x[n] W_N^{kn}, \quad k=0, 1, 2, \dots, N-1$$

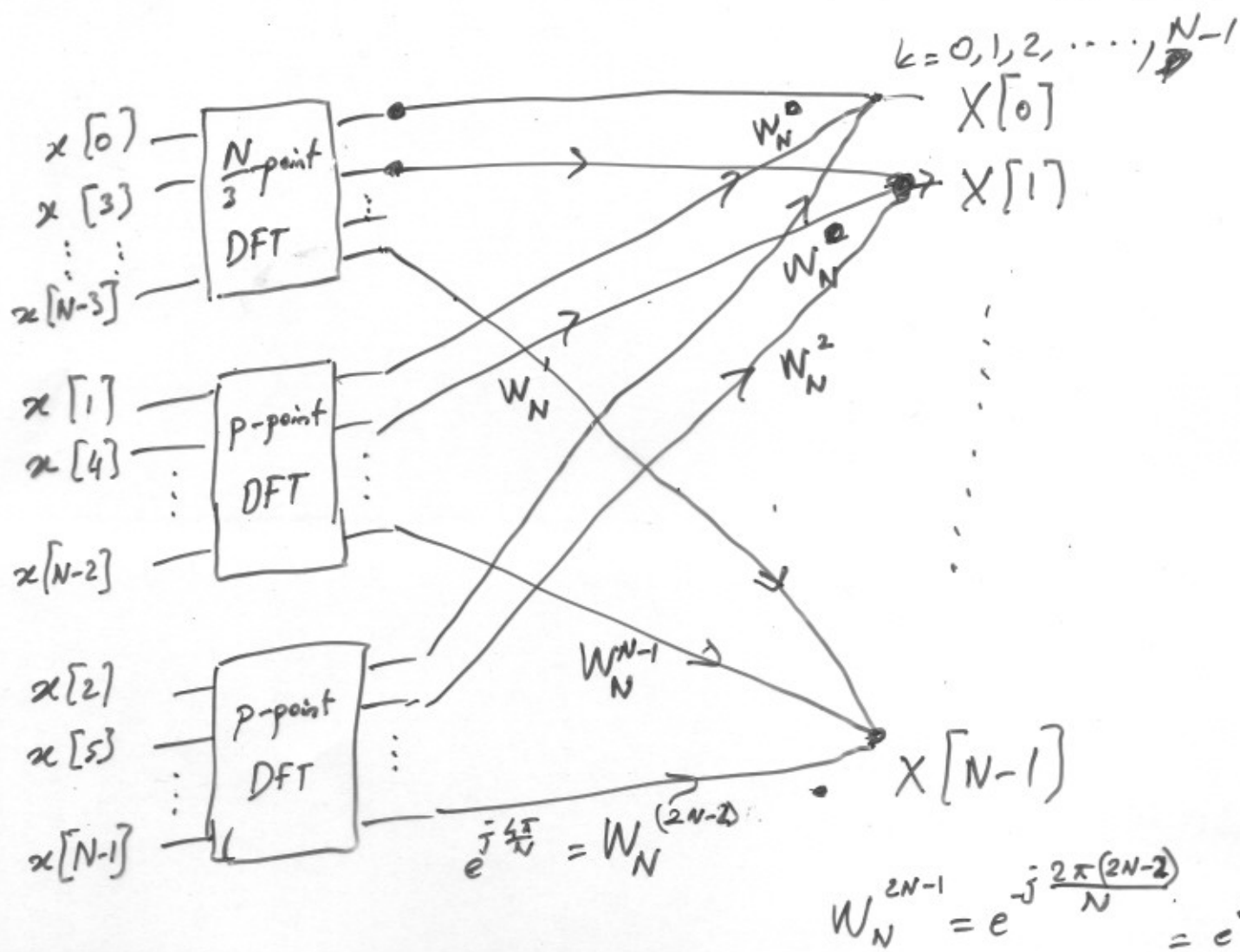
Decimation-in-time:

$$X[k] = \sum_{l=0}^{\frac{N}{3}-1} x[3l] e^{-j \frac{2\pi k}{N} 3l} + \sum_{l=0}^{\frac{N}{3}-1} x[3l+1] e^{-j \frac{2\pi k}{N} (3l+1)} + \sum_{l=0}^{\frac{N}{3}-1} x[3l+2] e^{-j \frac{2\pi k}{N} (3l+2)}$$

$(x[0], x[3], \dots)$ $(x[1], x[4], \dots)$ $(x[2], x[5], \dots)$

$$X[k] = \sum_{l=0}^{\frac{N}{3}-1} x[3l] e^{-j \frac{2\pi k}{(N/3)} l} + e^{-j \frac{2\pi k}{N} \frac{N}{3}} \sum_{l=0}^{\frac{N}{3}-1} x[3l+1] e^{-j \frac{2\pi k}{(N/3)} l} + e^{-j \frac{2\pi k}{N} 2 \frac{N}{3}} \sum_{l=0}^{\frac{N}{3}-1} x[3l+2] e^{-j \frac{2\pi k}{(N/3)} l}$$

$\frac{N}{3} - \text{DFT}$ $\frac{N}{3} = p - \text{point}$ $p - \text{point DFT's}$



Inverse D.F.T: $x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn}$

$$X[k] = \frac{1}{N} \sum_{n=0}^{N-1} x[n] e^{+j \frac{2\pi kn}{N}}, \quad n=0, 1, \dots, N-1$$

Proof:

Forward D.F.T:

$$\begin{bmatrix} X[0] \\ X[1] \\ \vdots \\ X[N-1] \end{bmatrix} = \begin{bmatrix} 1 & 1 & \dots & 1 \\ 1 & e^{j \frac{2\pi}{N}} & \dots & e^{j \frac{2\pi(N-1)}{N}} \\ \vdots & \vdots & \ddots & \vdots \\ 1 & e^{-j \frac{2\pi(N-1)}{N}} & \dots & e^{-j \frac{2\pi(N-1)^2}{N}} \end{bmatrix} \begin{bmatrix} x[0] \\ x[1] \\ \vdots \\ x[N-1] \end{bmatrix}$$

$\underbrace{\hspace{10em}}_{\underline{W}_N}$

(The last row $(1 \ e^{-j \frac{2\pi(N-1)}{N}} \ e^{-j \frac{2\pi(N-2)}{N}} \ \dots \ e^{-j \frac{2\pi}{N}})$)
 * \underline{W}_N is a symmetric matrix.
 * Rows of \underline{W}_N are orthogonal to each other.

$$W_N^{-1} = \frac{1}{N} W_N^H = \frac{1}{N} (W_N^*)^T = \frac{1}{N} W_N^*$$

$$\Rightarrow X[k] = \frac{1}{N} \sum_{n=0}^{N-1} x[n] e^{+j \frac{2\pi kn}{N}} \quad n=0, 1, \dots, N-1$$

Ex 11

$$x[n] = \begin{cases} 1, & n=0 \\ 0, & \text{o.w.} \end{cases}$$

$$\rightarrow X[k] = ?, \quad k=0, \dots, N-1$$

$$X[0] = \sum_{n=0}^{N-1} x[n] = 1$$

$$X[1] = \sum_n x[n] e^{-j \frac{2\pi kn}{N}} = 1 \cdot e^{-j \frac{2\pi 1 \cdot 0}{N}} + 0 + \dots + 0 = 1$$

$$X[N-1] = \sum_n x[n] e^{-j \frac{2\pi (N-1) \cdot n}{N}} = 1 \cdot e^{-j \frac{2\pi (N-1) \cdot 0}{N}} + 0 + \dots + 0 = 1$$

Inverse D.F.T. expression

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} \underbrace{1}_{X[k]} e^{+j \frac{2\pi kn}{N}} = \begin{cases} 1 & \text{for } n=0, N, 2N, \dots, 3N \\ 0 & \text{o.w. } (n=1, 2, \dots, N-1) \end{cases}$$

Fast Fourier Transform (FFT) Algorithm

DFT: $X[k] = \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi k}{N} n}$, $k=0, 1, \dots, N-1$

Direct implementation comp. cost: N^2 (complex multiplications) + N additions

* If $N = 1024 = 2^{10} \Rightarrow N^2 \approx 10^3 \times 10^3 = 10^6$ multiplications.

* FFT: $O(N \log_2 N)$ complex multiplications $\approx \frac{10^3}{2} \log_2 1024 \approx \frac{10^4}{2}$ real multiplications

Decimation-in-time FFT algorithm: ($W_N^k = e^{j \frac{2\pi k}{N}}$)
 [N is a power of 2].

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi k}{N} n} + \sum_{n=0}^{N-1} x[n] e^{-j \frac{2\pi k}{N} n}, k=0, 1, \dots, N-1$$

$$X[k] = \sum_{l=0}^{N/2-1} x[2l] e^{-j \frac{2\pi k}{N} 2l} + \sum_{l=0}^{N/2-1} x[2l+1] e^{-j \frac{2\pi k}{N} (2l+1)}$$

$$X[k] = \sum_{l=0}^{N/2-1} x[2l] (W_N^2)^{lk} + W_N^k \sum_{l=0}^{N/2-1} x[2l+1] e^{-j \frac{2\pi k}{N} 2l} (W_N^2)^{lk}$$

$$W_N^2 = W_{N/2} = e^{-j \frac{2\pi k}{N}} = e^{-j \frac{2\pi k}{N/2}} = W_{N/2}$$

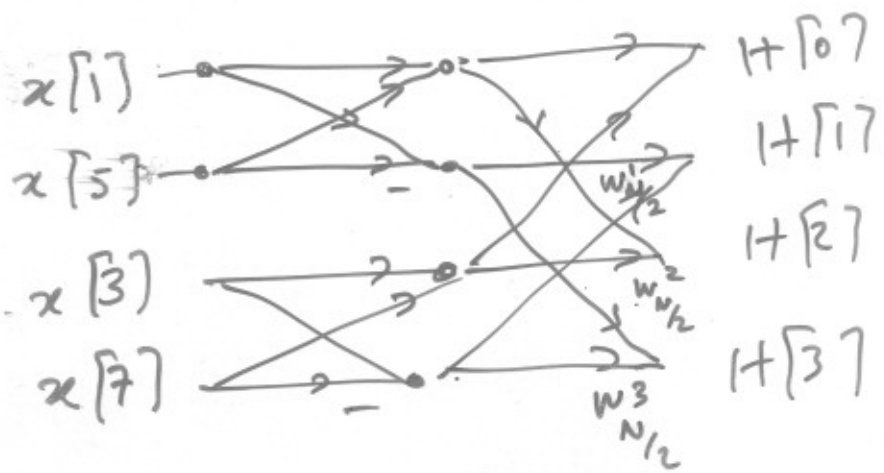
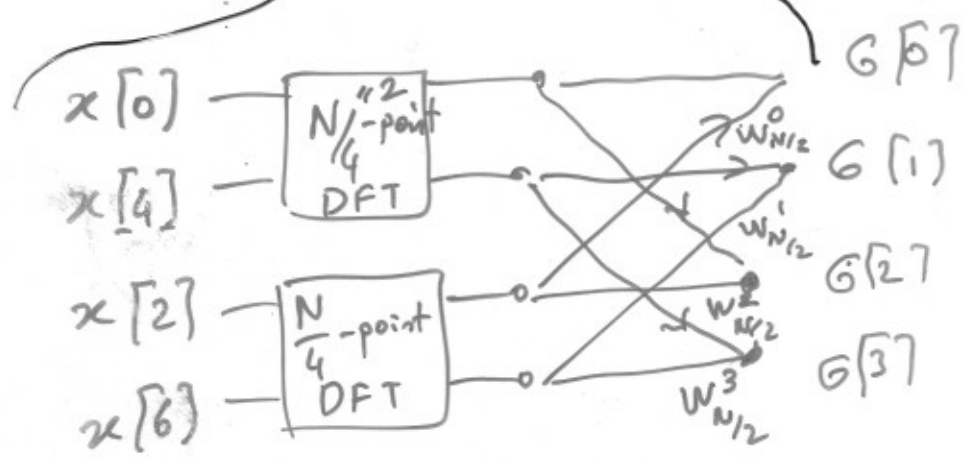
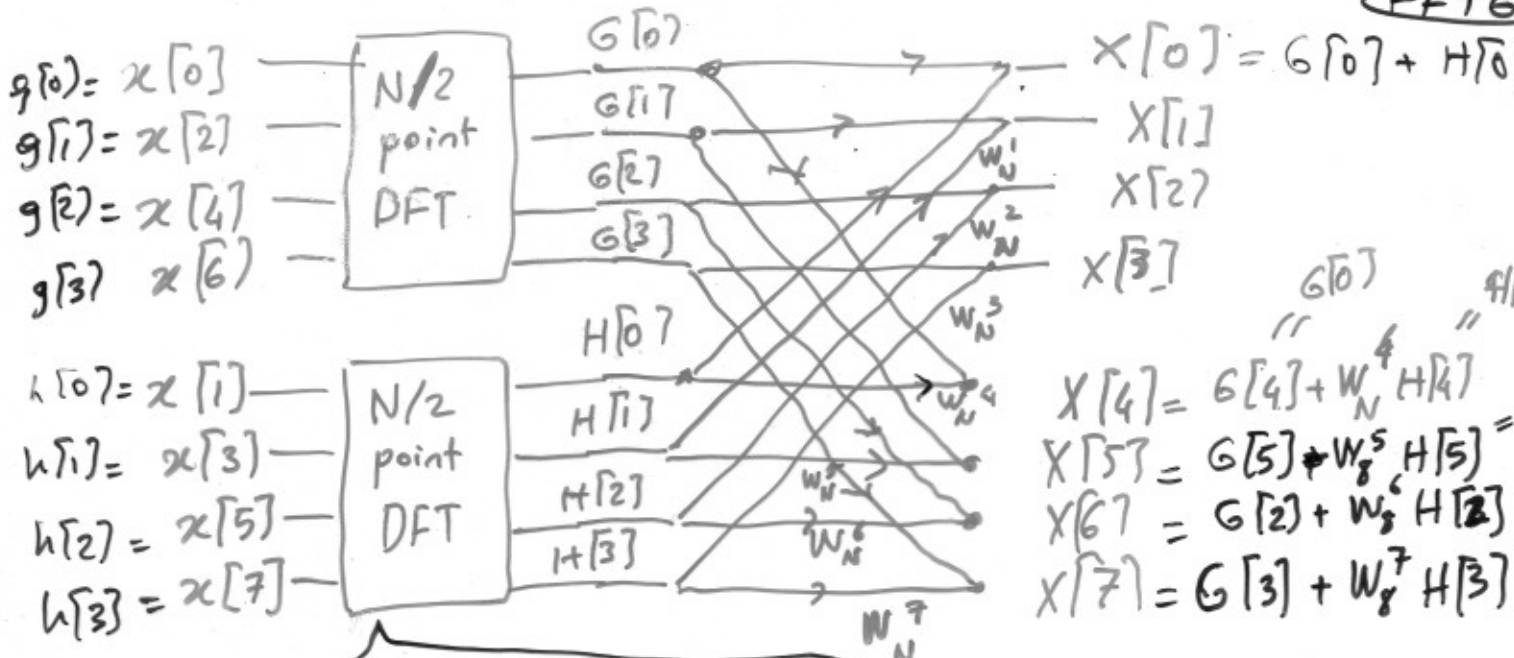
$$X[k] = \sum_{l=0}^{N/2-1} x[2l] W_{N/2}^{lk} + W_N^k \sum_{l=0}^{N/2-1} x[2l+1] W_{N/2}^{lk}$$

\swarrow $N/2$ point DFT of evens. \swarrow $N/2$ point DFT of odds

$$X[k] = G[k] + W_N^k H[k], k=0, 1, \dots, N-1 \quad (*)$$

Comp. cost study:
 $(N=1024=2^{10} \approx 10^3) \Rightarrow X[k]: \frac{(512)^2}{6[k]} + \frac{(512)^2}{H[k]} + 10^3 \approx 10^6 X[k]$
 periodic with $N/2 = 512 < 10^3$

$N=8$



Decimation-in-time 8-point DFT Flowchart.

$\neq \log_2 N$ stages. ($2^3 = 8 \Rightarrow 3$ stages)

* Input: bit reversed order \Rightarrow output: regular

$x[6] = 110$ - - reverse bits $x[3] = 011$